# Introduction to Bioinformatics for Computer Scientists 

## Lecture 6

## Plan for next lectures

- Today: Introduction to phylogenetics
- Lecture 7 (Alexis): Phylogenetic search algorithms


## The story so far

- Biological Terminology: RNA, DNA, genes, genomes, etc
- Pair-wise Sequence Alignment
- Sequence Comparison
- Genome Assembly
- Multiple Sequence Alignment


## The story so far

- Biological Terminology: RNA, DNA, genes, genomes, etc
- Pair-wise Sequence Alignment
- Sequence Comparison
- Genome Assembly
- Multiple Sequence Alignment
- Phylogenetic Inference


## A Taxonomy



First systematic classification of living beings by Aristotele 384-382 BC Some terms still in use today, e.g., classification of animals into Vertebrates versus Invertebrates


Wirbeltiere


## Taxonomy

- Group biological organisms (species) into groups with similar characteristics
- Define characteristics of groups at different hierarchy levels, e.g., animals > mammals > great apes
- Taxonomic ranks
- Domain $\rightarrow$ three domains of life
- Kingdom
- Phylum
- Class
- Order
- Family
- Genus
- Species


## A Phylogeny or Phylogenetic Tree



## A Phylogeny or Phylogenetic Tree



## Phylogeny

- An unrooted strictly binary tree
- Leafs are labeled by extant "übrig geblieben" (currently living) organisms represented by their DNA/Protein sequences
$\rightarrow$ we can also sequence ancient DNA, see, for instance, the neandertal genome: "The complete genome sequence of a Neanderthal from the Altai Mountains", Nature 2013
$\rightarrow$ depends on temperature, time, and other environmental conditions
$\rightarrow$ up to 300,000 years back, see http://www.pnas.org/content/110/39/15758.abstract
- Inner nodes represent hypothetical common ancestors
- Outgroup: one or more closely related, but different species $\rightarrow$ allows to root the tree


## Taxon

- Used to denote clades/subtrees in phylogenies or taxonomies
- A group of one or more species that form a biological unit
- As defined by taxonomists
$\rightarrow$ subject of controversial debates
$\rightarrow$ part of the culture/fuzziness of Biology
- In phylogenetics we often refer to a single leaf as taxon
$\rightarrow$ the plural of taxon is taxa
$\rightarrow$ we often say that a tree with $n$ leaves (sequences) has $n$ taxa


## Some more terminology

This phylogeny has a root!


B and C are a monophyletic group; they are sister species

## Some more terminology


( $\mathbf{A}, \mathbf{B}, \mathbf{C}$ ) is a monophyletic group; it is sister to ( $\mathbf{D}, \mathrm{E}$ )

## Some more terminology


$(A, B, C, D)$ is paraphyletic $\rightarrow E$ is excluded

## Some more terminology


$(A, D)$ is a polyphyletic group $\rightarrow$ their most recent common ancestor (MRCA) is excluded

## Some more terminology



Tree-based or patristic distance between two taxa:
Sum over branch lengths along the path in the tree, e.g.:

## Some more terminology



Tree-based or patristic distance between two taxa:
Sum over branch lengths along the path in the tree, e.g.:
A $\leftrightarrow \mathrm{B}: 0.2$

## Some more terminology



Tree-based or patristic distance between two taxa:
Sum over branch lengths along the path in the tree, e.g.:
A $\leftrightarrow$ B: 0.2
A $\leftrightarrow \mathbf{D}: 0.35$

## Tree Rooting

Outgroup species 1

Ingroup species 3
Outgroup species 2

Ingroup species 2
Ingroup species 1

## Tree Rooting



## Tree Rooting



## Tree Rooting



## Outgroup Choice



Ingroup species 4

Ingroup species 2

Ingroup species 3

Ingroup species 1
Clear signal

Close Outgroup

## Tree Inference



## Tree Inference



## Tree Inference



## Tree Inference



## Tree Inference

Taxon 1:ACGTTT
Taxon 2:ACGTT
Taxon 3:ACCCT

Taxon 4:AGGGTTT $\rightarrow$\begin{tabular}{c}
MSA <br>
Program

$\rightarrow$

Taxon 1:ACGTTT- <br>
Taxon 2:ACGTT- <br>
Taxon 3:ACCCT- <br>
Taxon 4:AGGGTTT

$\quad$

Tree inference <br>
program
\end{tabular}

## Tree Inference



## Tree Inference



Newick format with branch lengths:
(Taxon1:0.1,Taxon2:0.2,(Taxon3:0.15,Taxon4:0.15):0.3);

## Problems with Newick tree format

- Except for branch length values: no way to associate meta-data to branch lengths
- However, there is important meta-data, e.g., branch support: how well is a branch in the tree supported?
$\rightarrow$ ad hoc solution: represent branch support values as node meta-data!
$\rightarrow$ this causes problems


## Problems with Newick tree format



Branch support values represented as node meta-data can be assigned incorrectly to branches after re-rooting.

About $50 \%$ of the tools we checked had this Problem. For details see:
https://academic.oup.com/mbe/article/34/6/1535/3077051


Which representation is correct?

## A real example


a) original tree
b) re-rooted tree with shifted support values
c) re-rooted tree with correct support values

## Tree Shapes



## Tree Shapes



Non-ultrametric tree

## Tree Shapes



## Tree Shapes



## Dating Trees



## Dating Trees



## Dating Trees



## Dating Trees



We need a rooted \& ultrametric tree!
$\rightarrow$ rooting with outgroups
$\rightarrow$ ultrametricity with programs for divergence time estimation
$\rightarrow$ active research area
$\rightarrow$ most codes rely on the phylogenetic likelihood function and Bayesian Statistics (MCMC methods)

## Dating Trees



But how do we place the fossil?
$\rightarrow$ typically no DNA data available
Fossil placement:
$\rightarrow$ ad hoc using empirical knowledge
$\rightarrow$ computationally using morphological data

The input for a phylogenetic analysis need not be molecular data!

We can also use sequences of morphological traits ("Merkmale")!

## Remember that we deal with extant species!



## Morphological Traits

```
t1: 1000
t2: 0100
t3: 0010
T4: 0001
or:
t1: 0
t2: 1
t3: 2
t4: 3
```

t1: 1000
t2: 0100
t3: 0010
T4: 0001
or:
t1: 0
t2: 1
t3: 2
t4: 3
without hair removal)?


## Morphological Traits

What image best matches the extent of your natural brow line
t1: 1000
t2: 0100
t3: 0010
T4: 0001
or:


Traits need not be discrete, they can also be continuous, e.g., bone ratios

## Alignment-Free Tree Inference



## Alignment-Free Tree Inference



## How many unrooted 4-taxon trees exist?



## How many rooted 4-taxon trees exist?



## Tree Counts

- Unrooted binary trees
- 4 taxa $\rightarrow 3$ distinct trees
- A tree with $n$ taxa has $n-2$ inner nodes
- And 2n-3 branches
- Rooted binary trees
- 4 taxa $\rightarrow 3$ unrooted trees * 5 branches each (rooting points) $=15$ trees
- n-1 inner nodes
- $2 n-2$ branches


## The number of trees

3 taxa $=1$ tree

## The number of trees



4 taxa: 3 trees
u: \# trees of size 4-1 := 1
v : \# branches in a tree of size 4-1 := 3
Number of unrooted binary trees with 4 taxa: $u$ * $v=3$

## The number of trees



5 taxa: 15 trees
$\mathrm{u}=3$
$v=5$
Number of unrooted trees with 5 taxa: $3 * 5=15$

## The number of trees



## The number of trees explodes!



## Some Numbers

| Number of Organisms | Number of alternative Trees |
| :---: | :---: |
| 3 | 1 |
| 4 | 3 |
| 5 | 15 |
| 6 | 105 |
| 7 | 945 |
| 10 | 2.027 .025 |
| 15 | $7.905,853.580 .625$ |
| 20 | $2.21 * 10^{20}$ |
| 50 | $2.84 * 10^{76}$ |

Table 2.1: Number of possible trees for phylogenies with 3-50 organisms

## Equation for the number of unrooted trees

- Simple proof via induction

$$
\prod_{i=3}^{n}(2 i-5)
$$

- The number of rooted trees for $n$ taxa simply is the number of unrooted trees for $n+1$ taxa
- The additional $\left(n+1^{\text {th }}\right)$ taxon represents all possible rootings for all unrooted trees with $n$ taxa


## \# trees with 2000 tips

stamatak@exelixis:~/Desktop/GIT/TreeCounter\$ ./treeCounter -n 2000

GNU GPL tree number calculator released June 2011 by Alexandros Stamatakis




































 350341796875

## A side-note The treeCounter tool

- Evidently, the tree count can not be computed using normal integers
$\rightarrow$ we need an arbitrary precision library
$\rightarrow$ I used the GNU GMP (Multiple Precision Arithmetic) library
$\rightarrow$ treeCounter available as open-source code at https://github.com/stamatak
$\rightarrow$ Has anybody already used GNU GMP?


## Scoring Trees

- Now we know how many unrooted candidate trees there exist for $n$ taxa
- How do we chose among them?
$\rightarrow$ we need some scoring criterion $f()$ to evaluate them
$\rightarrow$ finding the optimal tree under most criteria is NP-Hard



## What can we do with Phylogenies?


reference phylogeny

## What can we do with Phylogenies?



## What can we do with Phylogenies?

Note that, this is similar to placing an outgroup into the tree!

reference phylogeny

## Diversification Rates



From: Charles C. Davis, Hanno Schaefer: "Plant Evolution: Pulses of Extinction and Speciation in Gymnosperm Diversity", Current Biology, 2011.

## Diversification Rates

- With former PostDoc Stephen Smith: "Understanding angiosperm diversification using small and large phylogenetic trees", American Journal of Botany 98 (3), 404-414, 2011.
- Largest tree of angiosperms computed to date
- 55,000 taxa



## Diversification Rates

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## Influenza Outbreaks

Host Taxa

- Galliformes

Anseriformes
Passeriformes - Charadriformes

Human
Columbidae
Artiodactyla
Accipitriformes
Ardeidae
Carnivora
O Corvidae
Arthropoda
Ambiguous


## And of course SARS-CoV-2

## Phylogenetic analysis of SARS-CoV-2 data is difficult

Benoit Morel ${ }^{+1}$, , Pierre Barbera ${ }^{*-1}$, Lucas Czech ${ }^{1}$, Ben Bettisworth ${ }^{1}$, Lukas Hübner ${ }^{12}$, Sarah Lutteropp ${ }^{1}$, Dora Serdarl ${ }^{17}$, Evangelia-Georgia Kostaki', loannis Mamais', Alexey M Kozlov', Pavios Pavilidis ${ }^{4}$, Dimittios Paraskevis', and Alexandros Stamatakis ${ }^{1,2}$



[^0]
## Snakebites

Australia has more poisonous snakes than any other continent, and many people die from snakebites each year. Developing effective antivenins is thus a high priority, but little is known about the venins of most species.
Phylogenetic analysis is helping with this task because venin properties correlate strongly with evolutionary relationships.
Although the red-bellied black snake looks very different from the king brown, it is actually closely related and can be treated with the same antivenin.
Conversely, the western brown looks very similar to the king brown, but it is only distantly related and thus responds best to different antivenin.
The phylogeny is also predictive: the recent demonstration that the poorlyknown barclick is closely related to the death adder (orange lineage) predicts that the former is also highly dangerous and might respond to widely-available death adder antivenin.


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## What can we do with phylogenetic trees?

- identifying unknown species
- divergence time estimates
- diversification rates
- viral outbreaks
- forensics $\rightarrow$ M.L. Metzker, D.P. Mindell, X.M. Liu, R.G. Ptak, R.A. Gibbs, D.M. Hillis: "Molecular evidence of HIV-1 transmission in a criminal case" PNAS: 99(22):14292-7, 2002.


## "Nothing in Biology makes sense, except in the light of evolution"

Why this increase in Phylogenetics papers? Advances in:
-Sequencing technology

- Hardware
- Methods \& Tools



## Building Trees

- We distinguish between
- Distance-based methods
$\rightarrow$ use MSA to compute a matrix of pair-wise distances
$\rightarrow$ build a tree using these distances
$\rightarrow$ Heuristics (essentially hierarchical clustering methods)
$\rightarrow$ Neighbor Joining: NJ
$\rightarrow$ Unweighted Pair Group Method with Arithmetic Mean: UPGMA
$\rightarrow$ least-squares method: explicit optimality criterion
- Character-based methods
$\rightarrow$ optimality criteria $f()$ operate directly on the MSA \& tree
$\rightarrow$ parsimony
$\rightarrow$ maximum likelihood
$\rightarrow$ Bayesian inference
$\rightarrow$ take the current tree topology \& MSA to calculate a score
$\rightarrow$ the score tells us how well the MSA data fits the tree


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Slow, but more accurate
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$\rightarrow$ Bayesian inference
Memory-intensive!
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$\rightarrow$ Unweighted Pair Group Method with Arith Lean: UPGMA
$\rightarrow$ least-squares method: explicit optimality

Less accurate, but faster

- Character-based methods
$\rightarrow$ optimality criteria $f($ ) operate direct

What could be the computational limitation here?
$\rightarrow$ parsimony
$\rightarrow$ maximum likelihood
$\rightarrow$ Bayesian inference
Memory-intensive!
$\rightarrow$ take the current tree topology \& MSA to calculate a score
$\rightarrow$ the score tells us how well the MSA data fits the tree

Slow, but more accurate

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## Out-of-core Algorithms

- Definition from Wikipedia:

Out-of-core or External memory algorithms are algorithms that are designed to process data that is too large to fit into a computer's main memory at one time. Such algorithms must be optimized to efficiently fetch and access data stored in slow bulk memory such as hard drive or tape drives.

- We do the data transfer RAM $\leftrightarrow$ disk explicitly from within the application code by using application-specific knowledge (e.g., about the data access patterns)
- This is to circumvent the paging procedure that would normally be initiated by the OS
- Out-of-core algorithms are typically much faster than the applicationagnostic paging procedure carried out by the OS
- For an example from phylogenetics see:

Fernando Izquierdo-Carrasco, Alexandros Stamatakis: "Computing the Phylogenetic Likelihood Function Out-of-Core", IEEE HICOMB 2011 workshop, Anchorage, USA, May 2011.

## NP-Hardness

- Because of the super-exponential increase in the number of possible trees for $n$ taxa ...
- all interesting criteria on trees are NP-hard:
- Least squares
- Parsimony $\rightarrow$ discrete criterion
- Likelihood $\rightarrow$ statistical criterion
- Bayesian $\rightarrow$ integrate likelihood over entire tree space


## Search Space



## Neighbor Joining $\rightarrow$ Principle



Given a kind of distance matrix $D_{i, j}$ where $i, j=1 . . .4$

## Neighbor Joining $\rightarrow$ Principle


min


Given a kind of distance matrix $D_{i, j}$ where $i, j=1 . . .4$ Find minimum and merge taxa

## Neighbor Joining $\rightarrow$ Principle



Given a kind of distance matrix $D_{i, j}$ where $i, j=1 \ldots 4$
Find minimum and merge taxa
Compute a new distance matrix of size $n-1=3$
Find minimum

## Neighbor Joining $\rightarrow$ Principle


min

Given a kind of distance matrix $D_{i, j}$ where $i, j=1 \ldots 4$
Find minimum and merge taxa
Compute a new distance matrix of size $n-1=3$
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## Neighbor Joining $\rightarrow$ Principle



Given a kind of distance matrix $D_{i, j}$ where $i, j=1 . . .4$
Find minimum and merge taxa
Compute a new distance matrix of size $n-1=3$
Find minimum and merge taxa
Etc.
Space complexity: $O\left(n^{2}\right)$
Time complexity: $O\left(n^{3}\right)$
Key question: how do we compute distance between $X$ and $A$ or $X$ and $B$ respectively
$\rightarrow$ for progressive alignment we may align the profile of $X$ with all remaining sequences

## Neighbor Joining Algorithm

- For each tip compute

$$
u_{i}=\Sigma_{j} D_{i j} /(n-2)
$$

$\rightarrow$ this is in principle the average distance to all other tips
$\rightarrow$ the denominator is $n-2$ instead of $n$, see below why

- Find the pair of tips, $(i, j)$ for which $D_{i j}-u_{i}-u_{j}$ is minimal
- Connect the tips (i,j) to build a new ancestral node $X$
- The branch lengths from the ancestral node $X$ to $i$ and $j$ are:

$$
\begin{aligned}
& b_{i}=0.5 D_{i j}+0.5\left(u_{i}-u_{j}\right) \\
& b_{j}=0.5 D_{i j}+0.5\left(u_{j}-u_{i}\right)
\end{aligned}
$$

- Update the distance matrix:
$\rightarrow$ Compute distance between the new node $X$ and each remaining tip as follows:

$$
D_{i j, k}=\left(D_{i k}+D_{j k}-D_{i j}\right) / 2
$$

- Replace tips $i$ and $j$ by the new node $X$ which is now treated as a tip
- Repeat until only two nodes remain
$\rightarrow$ connect the remaining two nodes with each other


## Neighbor Joining Algorithm

|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |
| B |  | - | 12 | 18 |
| C |  |  | - | 14 |
| D |  |  |  | - |

## Neighbor Joining Algorithm

|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |
| B |  | - | 12 | 18 |
| C |  |  | - | 14 |
| D |  |  |  | - |


| i | $u_{i}$ |
| :---: | :---: |
| A | $(17+21+27) / 2=32.5$ |
| B | $(17+12+18) / 2=23.5$ |
| C | $(21+12+14) / 2=23.5$ |
| D | $(27+18+14) / 2=29.5$ |

Distance matrix, usually denoted as $D$

## Neighbor Joining Algorithm

|  | A | B | C | D |
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| A | - | 17 | 21 | 27 |
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| D | $(27+18+14) / 2=29.5$ |


|  | A | B | C | D |
| :---: | :---: | :---: | :---: | :---: |
| A | - | -39 | -35 | -35 |
| B |  | - | -35 | -35 |
| 4 C |  |  | - | -39 |
| D |  |  |  | - |
| $\mathrm{D}_{\mathrm{ij}}-\mathrm{u}_{\mathrm{i}}-\mathrm{u}_{\mathrm{j}}$ |  |  |  |  |

Usually denoted as Q matrix

## Neighbor Joining Algorithm

|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
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| D |  |  |  | - |
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|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |
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| C |  |  | - | 14 |
| D |  |  |  | - |


| i | $u_{i}$ |
| :---: | :---: |
| A | $(17+21+27) / 2=32.5$ |
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|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | -39 | -35 | -35 |
| B |  | - | -35 | -35 |
| C |  |  | - | -39 |



D

$$
D_{i j}-u_{i}-u_{j}
$$

## Neighbor Joining Algorithm

|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |
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| B | $(17+12+18) / 2=23.5$ |
| C | $(21+12+14) / 2=23.5$ |
| D | $(27+18+14) / 2=29.5$ |


|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | -39 | -35 | -35 |
| B |  | - | -35 | -35 |
| C |  |  | - | -39 |



D

$$
\begin{aligned}
& b_{c}=0.5 \times 14+0.5 \times(23.5-29.5)=4 \\
& b_{D}=0.5 \times 14+0.5 \times(29.5-23.5)=10
\end{aligned}
$$

## Neighbor Joining Algorithm

|  | A | B | C | D | X |
| :--- | :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |  |
| B |  | - | 12 | 18 |  |
| C |  |  | - | 14 |  |
| D |  |  |  | - |  |
| X |  |  |  |  | - |



## Neighbor Joining Algorithm

|  | A | B | C | D | X |
| :--- | :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |  |
| B |  | - | 12 | 18 |  |
| C |  |  | - | 14 |  |
| D |  |  |  | - |  |
| X |  |  |  |  | - |

$$
\begin{aligned}
D_{X A} & =\left(D_{C A}+D_{D A}-D_{C D}\right) / 2 \\
& =(21+27-14) / 2 \\
& =17 \\
& \\
D_{\text {XB }} & =\left(D_{C B}+D_{D B}-D_{C D}\right) / 2 \\
& =(12+18-14) / 2 \\
& =8
\end{aligned}
$$



## Neighbor Joining Algorithm

|  | A | B | C | D | X |
| :--- | :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 | 17 |
| B |  | - | 12 | 18 | 8 |
| C |  |  | - | 14 |  |
| D |  |  |  | - |  |
| X |  |  |  |  | - |

$$
\begin{aligned}
D_{X A} & =\left(D_{C A}+D_{D A}-D_{C D}\right) / 2 \\
& =(21+27-14) / 2 \\
& =17 \\
D_{\text {XB }} & =\left(D_{C B}+D_{D B}-D_{C D}\right) / 2 \\
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\end{aligned}
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## Neighbor Joining Algorithm

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## Neighbor Joining Algorithm

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| :--- | :--- | :--- | :--- |
| A | - | 17 | 17 |
| B |  | - | 8 |
| X |  |  | - |


| i | $u_{i}$ |
| :---: | :---: |
| A | $(17+17) / 1=34$ |
| B | $(17+8) / 1=25$ |
| X | $(17+8) / 1=25$ |



## Neighbor Joining Algorithm

|  | A | B | X |
| :--- | :--- | :--- | :--- |
| A | - | 17 | 17 |
| B |  | - | 8 |
| X |  |  | - |


| i | $u_{i}$ |
| :---: | :---: |
| A | $(17+17) / 1=34$ |
| B | $(17+8) / 1=25$ |
| X | $(17+8) / 1=25$ |


|  | A | B | X |
| :--- | :--- | :--- | :--- |
| A | - | -42 | -28 |
| B |  | - | -28 |
| X |  |  | - |



$$
D_{i j}-u_{i}-u_{j}
$$

## Neighbor Joining Algorithm

|  | A | B | X |
| :--- | :--- | :--- | :--- |
| A | - | 17 | 17 |
| B |  | - | 8 |
| X |  |  | - |


| i | $u_{i}$ |
| :---: | :---: |
| A | $(17+17) / 1=34$ |
| B | $(17+8) / 1=25$ |
| X | $(17+8) / 1=25$ |


|  | A | B | X |
| :--- | :--- | :--- | :--- |
| A | - | -42 | -28 |
| B |  | - | -28 |
| X |  |  | - |



$$
D_{i j}-u_{i}-u_{j}
$$

## Neighbor Joining Algorithm

|  | A | B | X |
| :--- | :--- | :--- | :--- |
| A | - | 17 | 17 |
| B |  | - | 8 |
| X |  |  | - |


| i | $u_{i}$ |
| :---: | :---: |
| A | $(17+17) / 1=34$ |
| B | $(17+8) / 1=25$ |
| X | $(17+8) / 1=25$ |


|  | A | B | X |
| :--- | :--- | :--- | :--- |
| A | - | -42 | -28 |
| B |  | - | -28 |
| X |  |  | - |



$$
D_{i j}-u_{i}-u_{j}
$$

$$
\begin{aligned}
& b_{A}=0.5 \times 17+0.5 \times(34-25)=13 \\
& b_{D}=0.5 \times 17+0.5 \times(25-34)=4
\end{aligned}
$$

## Neighbor Joining Algorithm

|  | A | B | X | Y |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 17 |  |
| B |  | - | 8 |  |
| X |  |  | - |  |
| Y |  |  |  |  |



## Neighbor Joining Algorithm

|  | A | B | X | Y |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 17 |  |
| B |  | - | 8 |  |
| X |  |  | - | 4 |
| Y |  |  |  |  |

$$
\begin{aligned}
D_{Y X} & =\left(D_{A X}+D_{B X}-D_{A B}\right) / 2 \\
& =(17+8-17) / 2 \\
& =4
\end{aligned}
$$



## Neighbor Joining Algorithm

$$
\begin{array}{ccc} 
& X & Y \\
X & - & 4 \\
Y & & -
\end{array}
$$

$$
\begin{aligned}
D_{Y X} & =\left(D_{A X}+D_{B X}-D_{A B}\right) / 2 \\
& =(17+8-17) / 2 \\
& =4
\end{aligned}
$$



## Neighbor Joining Algorithm

$$
\begin{array}{ccc} 
& X & Y \\
X & - & 4 \\
Y & & -
\end{array}
$$

$$
\begin{aligned}
D_{Y X} & =\left(D_{A X}+D_{B X}-D_{A B}\right) / 2 \\
& =(17+8-17) / 2 \\
& =1
\end{aligned}
$$



## Neighbor Joining Algorithm

|  | A | B | C | D |
| :--- | :--- | :--- | :--- | :--- |
| A | - | 17 | 21 | 27 |
| B |  | - | 12 | 18 |
| C |  |  | - | 14 |
| D |  |  |  | - |




[^0]:    - A (103)

    ㄷ.. A. 1 (393)
    A. $2(60)$
    A. 3 (68)

    - A. 4 (9)
    A. 5 (28)
    A. $6(6)$
    $B(15)$

    B (15)
    B. $1(3068)$

    - B. 1 (3068)
    B. $2(446)$
    B. $2(446)$
    B. $3(70)$
    B. $3(70)$
    B. $4(100)$

    ㅁ. $\begin{gathered}\text { B. } 4(100) \\ \text { B. } 6(76)\end{gathered}$
    B. $6(76)$
    B. $9(5)$
    B. $9(5)$
    B. $10(1)$
    B. 10 (1)
    B. 11 (344)
    B. 11 (34
    B. 12 (2)
    B. 12 (2)
    B. 15 (6)
    B. 15 (6)
    B. 16 (23)
    B. 16 (23
    B. 17 (5)
    $8.17(5)$
    $8.18(2)$
    B. 18 (2)
    B. 21 (2)
    B. $21(2)$
    $-8.23(21)$
    $-\quad 8$
    B. $23(21$
    B.
    B.
    (5)
    B. 26 (8)

    - 8.27 (3)

