# Introduction to Bioinformatics for Computer Scientists 

## Lecture 8

## Outline

- Last time:
- How to Compute the Likelihood of a tree
- How to compute the Likelihood efficiently: Felsenstein Pruning Algorithm


## The Felsenstein Pruning Algorithm



Post order traversal

## Felsenstein Pruning



## Felsenstein Pruning



## Felsenstein Pruning



## Felsenstein Pruning



## Felsenstein Pruning



## Felsenstein Pruning



## Felsenstein Pruning

AND (left branch/right branch)

$$
\vec{L}_{A}^{(k)}(c)=\left(\sum_{S=A}^{T} P_{A S}\left(b_{i}\right) \vec{L}_{S}^{(i)}(c)\right)^{\prime}\left(\sum_{S=A}^{T} P_{A S}\left(b_{j}\right) \vec{L}_{S}^{(j)}(c)\right)
$$



## Felsenstein Pruning

OR (along left branch)

$$
\vec{L}_{A}^{(k)}(c)=\left(\sum_{S=A}^{T} P_{A S}\left(b_{i}\right) \vec{L}_{S}^{(i)}(c)\right)\left(\sum_{S=A}^{T} P_{A S}\left(b_{j}\right) \vec{L}_{S}^{(j)}(c)\right)
$$



## Felsenstein Pruning

OR (along right branch)

$$
\vec{L}_{A}^{(k)}(c)=\left(\sum_{S=A}^{T} P_{A S}\left(b_{i}\right) \vec{L}_{S}^{(i)}(c)\right)\left(\sum_{S=A}^{T} P_{A S}\left(b_{j}\right) \vec{L}_{S}^{(j)}(c)\right)
$$



## Felsenstein Pruning



## Felsenstein Pruning

Likelihood at the root: $L_{i}=\pi_{A} P(A)+\pi_{C} P(C)+\pi_{G} P(G)+\pi_{T} P(T)$


## An Excellent Tool to viualize and revise this concept

- https://phylanim.univ-lyon1.fr/LikelihoodTreeComputation

The computation of the likelihood of a site along a tree with four tips.
Let's compute the likelihood of observing site pattern $\{\mathrm{T}, \mathrm{T}, \mathrm{G}, \mathrm{T}\}$ at the four tips of a phylogeny


Computation

$$
\begin{aligned}
& L_{5}^{i}=\left(\sum_{j \in(A, C, G, T)} p_{i j}\left(v_{1}\right) \times L_{1}^{j}\right) \times\left(\sum_{k \in(A, C, G, T)} p_{i k}\left(v_{2}\right) \times L_{2}^{k}\right) \text { where } i \in(A, C, G, T) \\
& L_{5}^{G}=(0.032) \times\left(\left(P_{G A}\left(v_{2}\right) \times L_{2}^{A}\right)+\left(P_{G C}\left(v_{2}\right) \times L_{2}^{C}\right)+\left(P_{G G}\left(v_{2}\right) \times L_{2}^{G}\right)+\left(P_{G T}\left(v_{2} \times L_{2}^{T}\right)\right)\right. \\
& L_{5}^{G}=(0.032) \times\left(\left(P_{G A}(0.25) \times 0.0\right)+\left(P_{G C}(0.25) \times 0.0\right)+\left(P_{G G}(0.25) \times 0.0\right)+\left(P_{G T}(0.25) \times 1.0\right)\right) \\
& L_{5}^{G}=(0.032) \times\left((0.055 \times 0.0)+(0.039 \times 0.0)+\left(P_{G G}(0.25) \times 0.0\right)+\left(P_{G T}(0.25) \times 1.0\right)\right)
\end{aligned}
$$



## Why is time-reversibility important?

$$
L=\sum_{S_{4}=A}^{T} \pi_{S_{4}} \sum_{S_{3}=A}^{T} P_{S_{4} S_{3}}\left(b_{1}\right) L_{S_{3}}^{(3)} \sum_{S_{5}=A}^{T} P_{S_{4} S_{5}}\left(b_{4}\right) L_{S_{5}}^{(5)}
$$



## Why is time-reversibility important?

$$
L=L^{\prime}=\sum_{S_{4}=A}^{T} \pi_{S_{4}} \sum_{S_{3}=A}^{T} P_{S_{4} S_{3}}\left(b_{1}+x\right) L_{S_{3}}^{(3)} \sum_{S_{5}=A}^{T} P_{S_{4} S_{5}}\left(b_{4}-x\right) L_{S_{5}}^{(5)}
$$



## Why is time-reversibility important?

$$
L=L^{\prime}=\sum_{S_{4}=A}^{T} \pi_{S_{4}} \sum_{S_{3}=A}^{T} P_{S_{4} S_{3}}\left(b_{1}+x\right) L_{S_{3}}^{(3)} \sum_{S_{5}=A}^{T} P_{S_{4} S_{5}}\left(b_{4}-x\right) L_{S_{5}}^{(5)}
$$

$$
b_{1}{ }^{\prime}:=0 \quad b_{4}^{\prime}:=b_{1}+b_{4} \quad \square L_{S 5}
$$

## Why is time-reversibility important?



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- Parallel Likelihood computations


## What's in the black box $P_{i j}(t) ?$

Instantaneous rate matrix $R$ !


## What's in the black box $P_{i j}(t)$ ?

What about the probabilities of staying in the current state?
$\rightarrow$ they are given by the properties of continuous Markov chains! e.g., $\lambda_{A A}=-\left(\lambda_{A C}+\lambda_{A G}+\lambda_{A T}\right) \rightarrow$ remember from lecture on Markov models: rows in the $R$ matrix need to sum to 0


## What's in the black box $P_{i j}(t) ?$



## What's in the black box $P_{i j}(t) ?$

Diagonal values are given by the off-diagonal
values ( R matrix property)
$\lambda_{A A}=-\left(\lambda_{A C}+\lambda_{A G}+\lambda_{A T}\right)$


## What's in the black box $P_{i j}(t) ?$



Equilibrium frequency vector $\pi=\left(\Pi_{A}, \Pi_{C} \Pi_{G}, \Pi_{T}\right)$ where $\Pi_{A}+\Pi_{C}+\Pi_{C}+\Pi_{T}=1$

## The simple Jukes-Cantor model

|  | A | C | G | T |
| :---: | :---: | :---: | :---: | :---: |
| A | * | $\lambda$ | $\lambda$ | $\lambda$ |
| C |  | * | $\lambda$ | $\lambda$ |
| G |  |  | * | $\lambda$ |
| T |  |  |  | * |

$$
\Pi=(1 / 4,1 / 4,1 / 4,1 / 4)
$$

## The Felsenstein 81 model

|  | A | C | G | T |
| :---: | :---: | :---: | :---: | :---: |
| A | * | $\lambda$ | $\lambda$ | $\lambda$ |
| C |  | * | $\lambda$ | $\lambda$ |
| G |  |  | * | $\lambda$ |
| T |  |  |  | * |

$$
\Pi_{i} \neq \Pi_{j}
$$

## Kimura 2-parameter model 1980



## HKY85



$$
\Pi_{i} \neq \Pi_{j}
$$

## GTR 1986



$$
\Pi_{i} \neq \Pi_{j}
$$

## GTR 1986



$$
\Pi_{i} \neq \Pi_{j}
$$

## GTR 1986



Note that these are relative rates, their values only matter relative to each other, so we can set $\zeta:=1.0$ by default. Although the GTR model has 6 rates, it only has 5 free parameters!

$$
\Pi_{i} \neq \Pi_{j}
$$

## Model Hierarchy



## GTR 1986



This is a rate matrix, time reversibility would require $\boldsymbol{\pi} r_{i j}=\pi r_{j i}$

$$
\Pi_{i} \neq \Pi_{j}
$$

## GTR 1986



This is a rate matrix, time reversibility would require $\boldsymbol{n r}_{i j}=\pi r_{j i}$ Solution: introduce a $Q$ matrix $Q$ := $\operatorname{diag(п)~} R$


## GTR 1986



$$
\Pi_{i} \neq \Pi_{j}
$$ Solution: introduce a $Q$ matrix $Q$ := $\operatorname{diag(п)~} R$


Then, $n r_{i j}=\pi r_{j i j}$ holds

This is a rate matrix, time reversibility would require $\pi r_{i j}=\pi r_{j i}$

## So how do we compute $P(t)$ from $Q$ ?

- As we have seen in the lecture on Markov chains:
$P(t)=e^{Q t}=I+Q t+1 / 2!(Q t)^{2}+1 / 3!(Q t)^{3}+\ldots$
- but this is unfortunately a matrix exponential :-(
- I will spare you the details, but in general, e.g., for GTR we need to apply an Eigenvector/Eigenvalue decomposition of $Q$ to calculate:

$$
P(t)=U \exp \left(\operatorname{diag}\left(\lambda_{i}\right) t\right) U^{-1}
$$

Matrix and inverse matrix of eigenvectors of $Q$

## So how do we compute $P(t)$ from $Q$ ?

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$$
P(t)=U \exp \left(\operatorname{diag}\left(\lambda_{j}\right) t\right) U^{-1}
$$

$$
\triangle
$$

Diagonal matrix of eigenvalues of $Q$, here the exponential function $\exp ()$ is invoked on scalar values!

## Likelihood Calculations

- So far, we have only seen how to calculate a likelihood on a
- given, fixed tree topology
- with given fixed branch lengths
- and given, fixed remaining model parameters
- Computing the maximum likelihood score, is much more complicated as it requires

1. functions for optimizing continuous parameters
2. functions for searching the discrete space of trees

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## Maximum Likelihood



## Maximum Likelihood



## Maximum Likelihood

Length: m


Commonly denoted as $Q$ matrix: transition probs for time $d t$, for time $t: P(t)=e^{Q t}$

## Maximum Likelihood



## Maximum Likelihood



## Maximum Likelihood


virtual root: vr

## Maximum Likelihood




## Maximum Likelihood



Conditional Likelihood Vectors


## Maximum Likelihood




## Post-order Traversal

## virtual root

## Post-order Traversal

## virtual root

## Post-order Traversal

## virtual root

| AGCC |  |
| :---: | :---: |
|  |  |
|  | G C C |
| 1.0 | 0.00 .00 .0 |
| 0.0 | 0.01 .01 .0 |
| 0.0 | 1.00 .00 .0 |
|  | 0.00 .00 .0 |

## Post-order Traversal

## virtual root

## Post-order Traversal

## virtual root



## What happens when we compute this inner vector?

$$
\vec{L}_{A}^{(k)}(c)=\left(\sum_{S=A}^{T} P_{A S}\left(b_{i}\right) \vec{L}_{S}^{(i)}(c)\right)\left(\sum_{S=A}^{T} P_{A S}\left(b_{j}\right) \vec{L}_{S}^{(j)}(c)\right)
$$



Position c

## Post-order Traversal

## virtual root

## Post-order Traversal

virtual root


## Post-order Traversal

virtual root


## Post-order Traversal

## virtual root


:-)

## Post-order Traversal

## virtual root


:-)

## Post-order Traversal

## virtual root


:-)

## Post-order Traversal

## virtual root

:-)
i-)

:-)

## Post-order Traversal

Overall likelihood: sum over logarithms of per-site likelihoods
virtual root
:-)

## Post-order Traversal

$$
L=\sum_{S_{4}=A}^{T} \pi_{S_{4}} L_{S_{4}}
$$



## Maximum Likelihood



## optimize branch lengths

## Branch Length Optimization

starting branch

## Branch Length Optimization

## starting branch

## Essentially we place the virtual root into this branch here

## Branch Length Optimization

## starting branch



## Newton Rapshon

- We want to find the branch length $b$ that maximizes the likelihood $L(b)$ of the tree
- For this, we want to know where the first derivative of $L(b)$ is 0
- To achieve this numerically we use the Newton-Raphson procedure for root finding deploying the first and second derivative of the likelihood $L^{\prime}(b)$ and $L^{\prime \prime}(b)$
- Note that, the likelihood only depends on branch $b$, all other model parameters ( $Q$ matrix, base frequencies, tree topology) remain fixed


## Derivatives of $L(b)$

- To compute the derivatives of $L(b)$, we essentially need to be able to compute the derivatives of $P(b)$ since the rest is just sums and does not depend on $b$
- Recall
$P(b)=e^{Q b}=U e^{\wedge b} U^{-1}$
- thus
$(P(b))^{\prime}=U \wedge e^{\wedge b} U^{-1}$
- and
$(P(b))^{\prime \prime}=U \wedge^{2} e^{\wedge b} U^{-1}$
- In practice we compute the derivatives of the log likelihood $\log (L(b))$, but it is essentially the same (see next slide)


## Derivatives of $\log (L(b))$

- $1^{\text {st }}$ derivative: $L(b)^{\prime} / L(b)$
- $2^{\text {nd }}$ derivative: $\left(L(b) L(b)^{\prime \prime}-\left(L(b)^{\prime \prime}\right)^{2}\right) / L(b)^{2}$


## Newton Raphson



Get point $x_{2}$ by tangent at $x_{1}$ which is given by $2^{\text {nd }}$ derivative $L(b)$ "


## Newton Raphson



An animation


## Branch Length Optimization

starting branch

## Branch Length Optimization

starting branch

## Branch Length Optimization

starting branch

## Branch Length Optimization

starting branch


## Branch Length Optimization

starting branch


## Branch Length Optimization

starting branch


## Branch Length Optimization

## starting branch



## Maximum Likelihood



## Maximum Likelihood


optimize model parameters


## Maximum Likelihood


optimize model parameters
Seq 1 Seq 3

Methods used for model parameter optimization (other than branch lengths)

1. BFGS
2. Brent's method
3. Expectation maximization approaches

## Numerical Optimization Procedures

- See chapters 9 \& 10 of: Numerical Recipes in C - The Art of Scientific Computing


## Basic Operations Maximum Likelihood

- Compute Conditional Likelihood Vector at an inner node
- Compute Likelihood at Virtual Root
- Optimize a Branch Length for a given Branch
- Optimize all Branch Lengths
- Optimize other Model Parameters


## Basic Operations Maximum Likelihood

- Compute Conditional Likelihood Vector at an inner node
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The optimizers are the tricky routines!

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## Protein Substitution Models

- The GTR Q matrix for protein data has 189 free parameters instead of just 5 (DNA)
- Estimating 189 rate parameters is difficult, time-consuming, and may lead to overparameterizing the model
- Instead, empirical models such as JTT, LG, WAG, MTMAM, etc. are used
- The $Q$ matrices are obtained by jointly optimizing model parameters on a large collection of reference alignments
- The models differ with respect to:
- the amount of data used to obtain them
- the type of data on which the models have been optimized
- e.g., dedicated models for HIV, FLU, Mammals
- the numerical optimization methods used
- Examples of general models:
- LG: Le \& Gascuel: "An Improved General Amino Acid Replacement Matrix"
- WAG: Whelan \& Goldman: "A General Empirical Model of Protein Evolution Derived from Multiple Protein Families Using a Maximum-Likelihood Approach"


## Rate Heterogeneity among Sites



- Among-site rate heterogeneity
- Biological phenomenon
$\rightarrow$ different sites/columns evolve at different speeds
- Need to accommodate this in our models


## Г-Distribution



## Г-Distribution



## Discrete Г-Distribution



## An Abstract View of $\Gamma$



This is the integral of the likelihood we approximate via discretization

$$
\operatorname{LnL}(i)=\log \left(\frac{1}{\mathbf{\Delta}} *\left(L_{0}+L_{1}+L_{2}+L_{3}\right)\right)
$$

## An Abstract View of $\Gamma$



4 times higher memory consumption

## An Abstract View of $\Gamma$



4 times more floating point operations

## 「 Model of Rate Heterogeneity with 4 discrete rates




|  |  |  |  |
| :---: | :---: | :---: | :---: |
| 1 | 8 | U | 4 |
| mucmumumins | NI | W | UUSEETE- -2 |
|  | MY | M11 |  |
| N4 | ENHMNMMHILST |  |  |
| TACETISETITE | AtICutiecche | UACHMULDEC | DUEItCuT |
|  |  |  | Mindum |
| HIMSMMTM | MMMmgnum | WhI | Mry |
|  | सHसMSM1 | \$110 | श15 |
| nMmNMMMMNHM | MIMMRMIMMNAN | Hactuminder | DUEIIEU |
|  | mhycreater-6 |  | -uctities |

## Mixture Models

- The $\Gamma$ model of rate heterogeneity is a simple example of socalled mixture models
- From Wikipedia: "In statistics, a mixture model is a probabilistic model for representing the presence of subpopulations within an overall population, without requiring that an observed data set should identify the sub-population to which an individual observation belongs. Formally a mixture model corresponds to the mixture distribution that represents the probability distribution of observations in the overall population."
- The 「 model gives us 4 discrete evolutionary rates over which we integrate (add) the likelihood for each site, without assigning a specific rate to a specific site


## Mixture Models

- We can also imagine to integrate the likelihood over a set of
- distinct $Q$ matrices
- distinct base frequencies
- or combinations thereof
- The LG protein substitution model is an example thereof:
- It uses 4 distinct empirical $Q$ matrices and 4 distinct sets of base frequencies $\pi$ over which we integrate just like for the $\Gamma$ model


## An example


$\begin{array}{ll}\text { Travel time observation } & \ldots-\text { Single model_Weibull } \\ \text { Mixture model_GMM2 } & \text {-.- }\end{array}$

Taken from: "Measuring Service Reliability Using Automatic Vehicle Location Data" $\rightarrow$ bus service reliability

## Heterotachous Models

One GTR model for the entire tree

## Heterotachous Models



## Heterotachous Models



## What is a partitioned dataset?

Multi-gene or whole-genome alignment

## What is a partitioned dataset?



Multi-gene or whole-genome alignment

## What is a partitioned dataset?



Multi-gene or whole-genome alignment

## What is a partitioned dataset?

|  |  | Gene 1 | Gene 2 | Gene 3 |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  |  | Gene 4 |
|  |  |  |  |  |
| $a_{0}$ | $a_{1}$ | $a_{2}$ | $a_{3}$ | $a_{4}$ |
| GTR $_{0}$ | GTR $_{1}$ | GTR $_{2}$ | GTR $_{3}$ | GTR $_{4}$ |
|  |  |  |  |  |

## What is a partitioned dataset?



Joint branch length estimate

## What is a partitioned dataset?



## What is a partitioned dataset?



## What is a partitioned dataset?



## Models and Parameters

- If we add an additional parameter to a model, the likelihood will become better
- However, this does not mean anything, as
- We might be over-parameterizing
- The key question is if the more complex model yields a different tree topology
- So, how do we determine the best-fit model for a given dataset?


## Nested models

- A particular model is said to be nested within a more complex model only if constraining parameter values of the latter yields the former!
- So, the model can only be constrained in one direction to determine if its nested!
- If I need to constrain both models for which I intend to assess nesting, they are not nested.
- Example: The F81 (equal rates, unequal stationary frequencies) and K2P (2 distinct rates, equal stationary frequencies) models are not nested within each other.
$\rightarrow$ This is because fixing the parameter values of either model does not yield the other model
- However, they are both nested within GTR


## Model Testing

- If models are nested we can use a likelihood ratio test
- Model $A$ is nested in model $B$ if parameters in model $A$ are a subset of the parameters in model $B$
- For instance: the Jukes Cantor (JC) model is nested in the General Time Reversible (GTR) model of nucleotide substitution
- $L R=P(D \mid A) / P(D \mid B)=L(A) / L(B)$
- $\Delta=\ln \left(L R^{2}\right)=2(\ln (L(A))-\ln (L(B))$
- Compare $\Delta$ to $x^{2}$ distribution with $k_{A}-k_{B}$ degrees of freedom to determine if the $\Delta$ is significant or not
- The degrees of freedom difference is the difference in the number of free parameters in the models
- How many free parameters do the JC and GTR models have?


## Model Testing

- If models are nested we can use a likelihood ratio test
- Model $A$ is nested in model $B$ if parameters in model $A$ are a subset of the parameters in modn_o
- For instance: th We are only allowed to compare likelihoods on the same data D! (GTR) model of
- $L R=P(D \mid A) / P(D \mid B)$
- $\Delta=\ln \left(L R^{2}\right)=2(\ln (L(f)-\ln (L(B))$
- Compare $\Delta$ to $x^{2}$ distribution with $k_{A}-k_{B}$ degrees of freedom to determine if the $\Delta$ is significant or not
- The degrees of freedom difference is the difference in the number of free parameters in the models
- How many free parameters do the JC and GTR models have?
$\rightarrow$ JC: 0
$\rightarrow$ GTR: 8


## What if Models are not nested?

- One can use other criteria such as
- Akaike Information Criterion (AIC)
- Bayesian Information Criterion (BIC)
- I will spare you the details, but the basic idea always is:
- Compute likelihood of alternative models
- Penalize the more parameter-rich models


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## Data Structures for unrooted Trees

- Unrooted trees with dynamically changing virtual roots need a dedicated tree data structure
- Why can the virtual root positions change dynamically?
- If we apply a topological move (NNI, SPR, TBR) will we have to re-compute all conditional likelihood vectors?


## Memory Organization: Conditional Likelihood Vectors with an Unrooted View



## Memory Organization: Conditional Likelihood Vectors with a Rooted View



## Memory Organization: CLVs with a Rooted View

New Virtual Root


## Memory Organization: CLVs with a Rooted View



## Memory Organization: Ancestral Vectors with a Rooted View

New Virtual Root


## Memory Organization: Tip Vectors



## Memory Organization: Tip Vectors



## Optimization of Likelihood Calculations

- Use SSE3 \& AVX vector intrinsics
- Also: GPUs, FPGAs
- Special implementations (why?) for computing CLVs:



## Optimization of Likelihood Calculations

- Use SSE3 \& AVX vector intrinsics
- Also: GPUs, FPGAs
- Special implementations (why?) for computing CLVs:


A lot of entries will be zero here if $i$ and/or $j$ are tips
$\rightarrow$ simplify calculations

## Repeating Patterns

Identical values, two times pattern AG
A.... A....
G .... G ....

## Repeating Patterns

Detect identical patterns and omit second computation
A.... A....
G .... G ....

## Repeating Patterns

Also, shorten CLV $\rightarrow$ less memory required


## A.... A....

G .... G ....

## Repeating Patterns (Repeats)

Also, shorten CLV $\rightarrow$ less memory required

A.... A....

G .... G ....

## Floating Point Numbers

- Machine numbers are an imperfect mapping of the infinite real numbers to a finite number of machine values!



## Floating Point Arithmetics: The Root of All Evil

- Computational science mostly relies on floating-point intensive codes
- How do we verify these codes?
- We stand on shaky grounds
- Scientists using those codes assume that there are no bugs
- Double precision arithmetics required for certain applications
- Who knows what de-normalized floating point numbers are?
$\rightarrow$ Please have a look at:
J. Björndalen, O. Anshus: "Trusting floating point benchmarks-are your benchmarks really data-independent?" Applied Parallel Computing. State of the art in Scientific Computing 2010; pp 178-188, Springer. and at my micro-benchmark at:
https://github.com/stamatak/denormalizedFloatingPointNumbers


## Floating Point Arithmetics: The Root of All Evil

- Computational science mostly reli

Why is this relevant when e codes talking about Maximum Likelihood?

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## Post-order Traversal



## Post-order Traversal



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## Loop Level Parallelism

 virtual root
$P[i]=f(Q[i], R[i])$

## Loop Level Parallelism

virtual root

$P[i]=f(Q[i], R[i])$

## Loop Level Parallelism

 virtual root

## Loop Level Parallelism

 virtual root

## Loop Level Parallelism

virtual root


## Parallel Post-order Traversal

Only need to synchronize at the root $\rightarrow$ MPI_Reduce() to calculate: $\Sigma \log \left(l_{i}\right)$
virtual root


## Parallel Post-order Traversal

## Overall Score $\Delta$ <br> $\rightarrow+\square$

$\Sigma \log \left(\mathrm{l}_{\mathrm{i}}\right)$


## Classic Fork-Join with Busy-Wait



## Synchronizations in RAxML with Pthreads

- RAxML Pthreads for a run time of about 10 seconds on 16 cores/threads
- 404 taxa 7429 sites: 194,000 Barriers
- 1481 taxa 1241 sites: 739,000 Barriers
- A paper on performance of alternative PThreads barrier implementations:
S.A. Berger, A. Stamatakis: "Assessment of Barrier Implementions for Fine-Grain Parallel Regions on Current Multi-core Architectures", IEEE Cluster 2010.


# Classic Fork-Join with Busy-Wait (model optimization) 



# Classic Fork-Join with Busy-Wait (model optimization) 



# Problems start with partitioned datasets! 



## Parallel Performance Problems

- They all start with partitioned datasets!
- How do we distribute partitions to processors?
- How do we calculate parameter changes?
- How much time does our broadcast take?
- Goal: Keep all processors busy all the time
$\rightarrow$ minimize communication and synchronization!


## Example

## Blue Gene Red Gene

Sequence 1

Sequence 5

## Data Distribution

## Orangutan Gorilla Chimp Homo Sapiens AGGA TTTT



## Data Distribution

## Orangutan Gorilla Chimp Homo Sapiens <br> AACG TTTT AAGG TTT- A-GG TTTT AGGA TTTT



## Data Distribution

\section*{Orangutan Gorilla Chimp Homo Sapiens <br> | AACG | TTTT |
| :--- | :--- |
| AAGG | TTT- |
| A-GG | TTTT |
| AGGA | TTTT |}



## Data Distribution I



## Data Distribution I

## Orangutan Gorilla Chimp Homo Sapiens

## AACG TTTT AAGG TTT-A-GG TTTT AGGA TTTT

Works well when we have more partitions than processors:
May lead to load imbalance not all processors obtain equal number of sites! Not all partitions have equal length!


## Data Distribution II

## Orangutan Gorilla Chimp Homo Sapiens

## AACG TTTT <br> AAGG TTT- <br> A-GG TTTT AGGA TTTT

Works well when we have more processors than partitions:
However we will need to compute: $P(t)=e Q t$ for each partition at each processor!


## Data Distribution II

## Orangutan Gorilla Chimp Homo Sapiens



Works well when we have more processors than partitions:
However we will need to compute: $P(t)=e Q t$ for each partition at each processor!


## Data Distribution II

## Orangutan Gorilla

Performance impact depends on number of states in data/dimension of $Q$


## Data Distribution II

## Orangutan Gorilla

How do we distribute partitions to processors?


## Load Balance I

| G0 | G1 | G2 | G3 |
| :--- | :--- | :--- | :--- |



## Load Balance I



## Load Balance I



## Load Balance I

- The multiprocessor job scheduling problem in phylogenetics
- Problem when \#partitions >> \#cores
- Tested per-site (cyclic/modulo) data distribution versus per partition data distribution
- We used the Longest Processing Time (LPT) heuristics for assigning partitions to processors
- 25 taxa, 220,000 sites, 100 genes
- GAMMA model
naïve: 613 secs

LPT: 550 secs

- CAT model
naïve: 298 secs
LPT: 127 secs
- Larger protein dataset under $\Gamma$ model of rate heterogeneity: 10-fold performance improvement!
J. Zhang, A. Stamatakis: "The Multi-Processor Scheduling Problem in Phylogenetics", 11th IEEE HICOMB workshop (in conjunction with IPDPS 2012).


## LPT heuristics for multi-processor scheduling

- Sort jobs (partitions) by processing length (partition length) in decreasing order
- Remove a job (partition) from the sorted list and assign it to the processor with the earliest end time (the smallest sum of partition lengths/number of sites)
- Repeat until the sorted list is empty
- Upper bound: $4 / 3-1 /(3 p)$ * OPT, where $p$ is the number of processors
- Graham, R. L.: "Bounds on Multiprocessing Timing Anomalies". SIAM Journal on Applied Mathematics 17 (2): 416-429, 1969.
- Remark: LPT works surprisingly well (see our paper on the phylogenetic problem where we also tested other heuristics)


## Partitioned Branch Lengths \& other parameters



## Load-Balance II



## Synchronization Points

- Assume 10 branches
- Each branch requires 10 Newton-Raphson Iterations
- Each NR Iteration requires a synchronization via a reduction operation
- One branch/partition at a time: 100 sync. points, less work (only one partition) per sync. point
- All branches concurrently: 10 sync. points, more work per sync. point
- Branches will need distinct number of operations
- Add convergence state $\rightarrow$ bit vector


# Synchronization Points 

$$
\begin{aligned}
& \text { Org1 AC GT } \\
& \text { Org2 AC TT }
\end{aligned}
$$

# Synchronization Points 

$$
\begin{array}{ll|l}
\hline \text { Org1 } & \text { AC } & \text { GT } \\
\text { Org2 } & \text { AC } & \text { TT }
\end{array}
$$

# Synchronization Points 

$$
\begin{array}{ll|l}
\text { Org1 } & \mathrm{AC} & \mathrm{GT} \\
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\end{array}
$$

# Synchronization Points 

$$
\begin{array}{ll|l}
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\hline
\end{array}
$$

# Synchronization Points 

$$
\begin{array}{ll|l}
\text { Org1 } & \mathrm{AC} & \mathrm{GT} \\
\text { Org2 } & \mathrm{AC} & \mathrm{TT}
\end{array}
$$

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# Synchronization Points 

$$
\begin{array}{ll|l}
\hline \text { Org1 } & \text { AC } & \text { GT } \\
\text { Org2 } & \text { AC } & \text { TT } \\
\hline
\end{array}
$$

In this example: 4 instead of 7 sync points!

## Load Balance II


A. Stamatakis, M. Ott: "Load Balance in the Phylogenetic Likelihood Kernel". Proceedings of ICPP 2009, Vienna, Austria, September 2009.

## Classic Fork-Join with



## Alternative MPI parallelization



